# A lightweight zerocopy notification mechanism

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# Agenda

- Background
- Problem Statement
- Design of Solution
- Implementation
- Evaluation

# Background - history of TCP zerocopy in Linux: TX

Linux had sendfile() support since early days.

This is using ops->sendpage() or ops->sendpage\_locked() and available to splice() users, like sendfile() and vmsplice()

In linux-4.14, Willem de Bruijn added MSG\_ZEROCOPY support to sendmsg() system call, along with completions sent to the socket error queue.

sendmsg(MSG\_ZEROCOPY) is slightly more efficient since it does not have to lock the socket for every page, unlike tcp\_sendpage().

## Background - history of TCP zerocopy in Linux: TX

The sendmsg zerocopy needs hardware to support standard SG support and TX checksum offloading. There is no requirements on how memory blocks need to be sized/aligned.

Almost all modern NIC support these (and more)

## Background - MSG\_ZEROCOPY notification

```
// Socket Setup
setsockopt(fd, SOL_SOCKET, SO_ZEROCOPY, &one, sizeof(one))
// Tranmission
ret = send(fd, buf, sizeof(buf), MSG_ZEROCOPY);
// Notification Reception
pfd.fd = fd;
pfd.events = 0;
if (poll(\&pfd, 1, -1) != 1 || pfd.revents \& POLLERR == 0)
        error(1, errno, "poll");
ret = recvmsg(fd, &msg, MSG_ERRQUEUE);
read_notification(msg);
```

# Background - MSG\_ZEROCOPY notification

```
read_notification(struct msghdr *msg, ...) {
        struct sock_extended_err *serr;
        struct cmsghdr *cm;
        cm = CMSG_FIRSTHDR(msg);
        if (cm->cmsg_level != SOL_IP &&
                cm->cmsg_type != IP_RECVERR)
                        error(1, 0, "cmsg");
        serr = (void *) CMSG_DATA(cm);
        if (serr->ee_errno != 0 ||
                serr->ee_origin != S0_EE_ORIGIN_ZEROCOPY)
                        error(1, 0, "serr");
        printf("completed: %u..%u\n", serr->ee_info, serr->ee_data);
```

#### **Problem Statement**

```
// Socket Setup
setsockopt(fd, SOL_SOCKET, SO_ZEROCOPY, &one, sizeof(one))
// Tranmission
ret = send(fd, buf, sizeof(buf), MSG_ZEROCOPY);
// Notification Reception
pfd.fd = fd;
pfd.events = 0;
if (poll(\&pfd, 1, -1) != 1 || pfd.revents \& POLLERR == 0)
        error(1, errno, "poll");
                                                               Extra overhead of
                                                               system calls
ret = recvmsg(fd, &msg, MSG_ERRQUEUE);
read_notification(msg);
```

#### **Problem Statement**

```
// Socket Setup
setsockopt(fd, SOL_SOCKET, SO_ZEROCOPY, &one, sizeof(one))

// Tranmission
ret = send(fd, buf, sizeof(buf), MSG_ZEROCOPY);

// Notification Reception
```

How to return some information back to the user upon returning of sendmsg without introducing extra system calls?

We already have similar mechanism in recvmsg.

## msg\_control in recvmsg

```
struct {
       struct cmsghdr cm;
        char control[512];
} control;
msg.msg_control = &control;
msg.msg_controllen = sizeof(control);
res = recvmsg(sock, &msg, recvmsg_flags|MSG_DONTWAIT);
if (res >= 0)
   printpacket(&msg, res, data, sock, ...);
```

## msg\_control in recvmsg

```
msg->msg_control
                   cmsg_hdr
                                   data
                                               cmsg_hdr
                                                              data
                                                                          cmsg_hdr
                                                                                          data
  static void printpacket(struct msghdr *msg, ...) {
      struct cmsghdr *cmsg;
      for (cmsg = CMSG_FIRSTHDR(msg);
          cmsq;
          cmsg = CMSG_NXTHDR(msg, cmsg)) {
          switch (cmsg->cmsg_type) {
              case SO TIMESTAMP: // print info
              case S0_TIMESTAMPNS: // print info
              case SO_TIMESTAMPING: // print info
```

# How about msg\_control in sendmsg?

```
msq.msq control = control;
msg.msg_controllen = sizeof(control);
cmsg = CMSG_FIRSTHDR(&msg);
cmsg->cmsg_level = SOL_SOCKET;
cmsg->cmsg_type = S0_TIMESTAMPING;
cmsg->cmsg_len = CMSG_LEN(sizeof(uint32_t));
val = sendmsg(fd, &msg, 0);
if (!cfg_busy_poll) {
    if (cfg_use_epoll)
        epoll(epfd);
    else
        __poll(fd);
while (!recv errmsg(fd)) {}
```

## Design - support msg\_control copy back to user in sendmsg

```
msg.msg_control = pointer;
msg.msg_controllen = CMSG_SPACE(size_of_placeholder))
cmsg = CMSG_FIRSTHDR(&msg);
cmsg->cmsg_level = SOL_SOCKET;
cmsg->cmsg_type = SCM_ZC_NOTIFICATION;
cmsg->cmsg_len = CMSG_LEN(size_of_placeholder);
ret = sendmsg(fd, &msg, 0);
if (ret >= 0)
       print_zc_info(&msg, ...)
```

# Design - support msg\_control copy back to user in sendmsg

```
print_zc_info(struct msghdr *msg, ...) {
        struct cmsghdr *cmsg;
        for (cmsg = CMSG_FIRSTHDR(msg);
             cmsg;
             cmsg = CMSG_NXTHDR(msg, cmsg)) {
                if (cm->cmsg_level != SOL_SOCKET &&
                    cm->cmsg_type != SCM_ZC_NOTIFICATION)
                        // deal with the zc notificatoin
```

## Design - compatibility

```
msg.msg_control = pointer;
msg.msg_controllen = CMSG_SPACE(...)
cmsg = CMSG_FIRSTHDR(&msg);
cmsg->cmsg_level = SOL_SOCKET;
cmsq->cmsq_type = SCM_ZC_NOTIFICATION;
cmsg->cmsg_len = CMSG_LEN(size_of_placeholder);
cmsg = CMSG_NXTHDR(\&msg);
                                                  Compatible with the original usage
cmsg->cmsg_level = SOL_SOCKET;
                                                  of msg control in sendmsg
cmsq->cmsq type = SO TIMESTAMPING;
cmsg->cmsg_len = CMSG_LEN(sizeof(uint32_t));
ret = sendmsg(fd, &msg, 0);
if (ret >= 0)
       print_zc_info(&msg, ...)
// Other logic related to SO_TIMESTAMPING
```

## Design - trailing notifications

```
while (...) {
   msg.msg_control = control;
   msg.msg_controllen = sizeof(control);
    cmsg = CMSG_FIRSTHDR(&msg);
    cmsg->cmsg_level = SOL_SOCKET;
    cmsg->cmsg_type = SCM_ZC_NOTIFICATION
    cmsg->cmsg_len = CMSG_LEN(sizeof(placeholder));
    ret = sendmsg(fd, &msg, 0);
    if (ret >= 0)
                                                     The zc notification of the last
        print_zc_info(&msg, res, data, sock, ...);
                                                     several sendmsgs might be empty
```

## Design - user interface - compatible with the original method

```
while (...) {
    msq.msq control = control;
    msq.msq controllen = sizeof(control);
    cmsg = CMSG_FIRSTHDR(&msg);
    cmsq->cmsq level = SOL SOCKET;
    cmsg->cmsg_type = SCM_ZC_NOTIFICATION
    cmsq->cmsq len = CMSG LEN(sizeof(placeholder));
    ret = sendmsg(fd, &msg, 0);
    if (ret >= 0)
        print_zc_info(&msg, res, data, sock, ...);
end:
// Trailing notification reception
pfd.fd = fd;
pfd.events = 0:
if (poll(&pfd, 1, -1) != 1 || pfd.revents & POLLERR == 0)
        error(1, errno, "poll");
ret = recvmsg(fd, &msg, MSG_ERRQUEUE);
// Notification parsing
read_notification(msg);
```

## Current msg\_control logic in sendmsg

```
struct msghdr {
        /*
        * Ancillary data. msg_control_user is the user buffer used for the
        * recv* side when msg_control_is_user is set, msg_control is the kernel
        * buffer used for all other cases.
       union {
               void
                               *msg_control;
               void user
                               *msg_control_user;
        };
        bool
                       msg_control_is_user : 1;
        . . .
```

## Current msg\_control logic in sendmsg

```
static int ____sys_sendmsg(...)
       unsigned char ctl[sizeof(struct cmsqhdr) + 20]
                                aligned(sizeof( kernel size t));
       if (ctl_len > sizeof(ctl)) {
               ctl_buf = sock_kmalloc(sock->sk, ctl_len, GFP_KERNEL);
               if (ctl buf == NULL)
                        goto out;
                if (copy_from_user(ctl_buf, msg_sys->msg_control_user, ctl len))
                        goto out freectl;
               msq sys->msq control = ctl buf;
               msq_sys->msg_control_is_user = false;
```

User passed in msg\_control\_user address is overwritten by a kernel buffer pointer. For the convenience of further access in the kernel.

## Implementation - A generic msg\_control copy back framework

```
static int ____sys_sendmsg(...)
       unsigned char ctl[sizeof(struct cmsqhdr) + 20]
                                aligned(sizeof( kernel size t));
        if (msg && msg_sys->msg_control_copy_to_user && err >= 0) {
                ssize_t len = err;
                err = sendmsg_copy_cmsg_to_user(msg_sys, msg);
                if (!err)
                        err = len;
out_freectl:
        if (ctl_buf != ctl)
                sock_kfree_s(sock->sk, ctl_buf, ctl_len);
out:
        return err;
```

## Implementation - A generic msg\_control copy back framework

```
static int sendmsg copy cmsg to user(struct msghdr *msg sys, struct user msghdr user *umsg)
{
    struct msqhdr msq user = *msq sys;
    msq user.msq control is user = true;
    msg_user.msg_control_user = umsg->msg_control;
    for_each_cmsghdr(cmsg, msg_sys) {
        if (!CMSG_OK(msg_sys, cmsg))
                break:
        if (cmsq copy to user(cmsq))
                put_cmsg(&msg_user, cmsg->cmsg_level, cmsg->cmsg_type,
                         cmsq->cmsq len - sizeof(*cmsq), CMSG DATA(cmsq));
```

## Implementation - A generic msg\_control copy back framework

```
static int sendmsg copy cmsg to user(struct msghdr *msg sys, struct user msghdr user *umsg)
   struct msqhdr msq user = *msq sys;
   msq user.msq control is user = true;
   msg_user.msg_control_user = umsg->msg_control;
   for_each_cmsghdr(cmsg, msg_sys) {
       if (!CMSG_OK(msg_sys, cmsg))
               break:
       if (cmsg_copy_to_user(cmsg))
               put_cmsg(&msg_user, cmsg->cmsg_level, cmsg->cmsg_type,
                        cmsg->cmsg_len - sizeof(*cmsg), CMSG_DATA(cmsg));
              put cmsg is used here to handle compat cases
```

#### Implementation - how to make use of the framework in zerocopy?

```
static inline bool cmsg_copy_to_user(struct cmsghdr *__cmsg) {
    return __cmsg->cmsg_type == SCM_ZC_NOTIFICATION;
int __sock_cmsg_send(struct sock *sk, struct cmsghdr *cmsg, ..) {
    switch (cmsg->cmsg_type) {
        case SCM_ZC_NOTIFICATION:
        // Get the information from MSG_ERRQUEUE
        // Populate the kernel buffer cmsg with the information
        msg->msg_control_copy_to_user = true;
```

#### Implementation - overall

```
static int ____sys_sendmsg(...)
        unsigned char ctl[sizeof(struct cmsghdr) + 20]
                                aligned(sizeof( kernel size t));
        msg_sys->msg_control_copy_to_user = false;
        /* In __sock_cmsg_send, if a cmsg needs to be copied back
        * to the user, handler function can update the kernel buffer
        * directly and set msg control copy to user to true.
        */
        err = __sock_sendmsg(sock, msg_sys);
        if (msg && msg_sys->msg_control_copy_to_user && err >= 0) {
                ssize t len = err;
                err = sendmsq_copy_cmsq_to_user(msq_sys, msq);
                if (!err)
                        err = len;
```

## Evaluation - msg\_zerocopy selftest

```
do {
    sends_since_notify++;
    do_sendmsg(fd, &msg, cfg_zerocopy, domain);

if (sends_since_notify >= cfg_notification_limit) {
        do_recv_completions();
        sends_since_notify = 0;
    }
} while (gettimeofday_ms() < tstop);</pre>
```

## Evaluation - throughput performance - notification interval = 1

Test Type / Protocol	TCP v4	TCP v6	UDP v4	UDP v6
ZCopy (MB)	7523	7706	7489	7304
New ZCopy (MB)	8834	8993	9053	9228
New ZCopy / ZCopy	117.42%	116.70%	120.88%	126.34%

## Evaluation - throughput performance - notification interval = 32

Test Type / Protocol	TCP v4	TCP v6	UDP v4	UDP v6
ZCopy (MB)	8842	8735	10072	9380
New ZCopy (MB)	9366	9477	10108	9385
New ZCopy / ZCopy	106.00%	108.28%	100.31%	100.01%

#### Evaluation - overhead introduced

```
static int sys sendmsq(...)
       unsigned char ctl[sizeof(struct cmsqhdr) + 20]
                               aligned(sizeof( kernel size t));
       msg sys->msg control copy to user = false;
       /* In __sock_cmsg_send, if a cmsg needs to be copied back
        * to the user, handler function can update the kernel buffer
        * directly and set msg_control_copy_to_user to true.
        */
       err = __sock_sendmsg(sock, msg_sys);
       if (msg && msg sys->msg control copy to user && err >= 0) {
               ssize_t len = err;
               err = sendmsg_copy_cmsg_to_user(msg_sys, msg);
               if (!err)
                       err = len;
                In the hot path, a minor cost is added to every
                other send calls which do not use this feature
```

## Next step - other possible use cases - timestamp

```
// Socket Setup
sock_opt = SOF_TIMESTAMPING_SOFTWARE | SOF_TIMESTAMPING_OPT_CMSG | SOF_TIMESTAMPING_OPT_ID
setsockopt(fd, SOL_SOCKET, SO_TIMESTAMPING, (char *)&sock_opt, sizeof(sock_opt))

// Tranmission
ret = send(fd, &msg, 0);

// Timestamp Reception
/* poll + recvmsg */
read_timestamps(msg);
```

## Next step - other possible use cases - Homa

```
int homa_send(...)
        struct homa_sendmsg_args args;
        args.id = 0;
        args.completion_cookie = completion_cookie;
        hdr.msg_control = &args;
        hdr.msg_controllen = 0;
        result = sendmsg(sockfd, &hdr, 0);
        // upon returning args.id needs to be updated
        if ((result >= 0) && (id != NULL))
                *id = args.id;
        return result;
```

## Next step - other possible use cases - Homa

```
int homa_send(...)
{
         hdr.msg_control = &args;
         hdr.msg_controllen = 0;
         result = sendmsg(sockfd, &hdr, 0);
         ...
}
```

# Summary

 A lightweight zerocopy notification mechanism to save the overhead of extra system calls.

- A generic msg\_control copy back framework in sendmsg, potentially apply to any other use cases where we need to return info back to user space.

#### **Patchset**

https://lore.kernel.org/all/20240708210405.870930-1-zijianzhang@bytedance.com/

Thanks for the code reviewing and suggestions by Willem de Bruijn and Xiaochun Lu!

Any comments or questions?

#### **Patchset**

https://lore.kernel.org/all/20240708210405.870930-1-zijianzhang@bytedance.com/

Thanks for the code reviewing and suggestions by Willem de Bruijn and Xiaochun Lu!

Open Question: Is this feature worth the minor cost in the sendmsg hot path?

## Appendix - design history

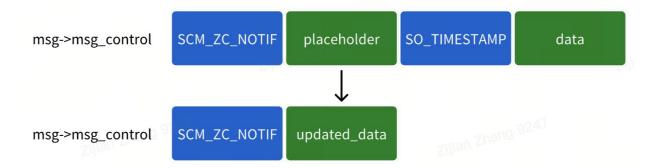
```
int __sock_cmsg_send(struct sock *sk, ...) {
    if (in_compat_syscall())
        usr_addr = compat_ptr(*(compat_uptr_t *)CMSG_DATA(cmsg));
   else
        usr_addr = (void __user *)*(void **)CMSG_DATA(cmsg);
    if (!access ok(usr addr, cmsg data len))
        return -EFAULT;
   // Retrieve zc notifications, and copy to a kernel buffer
    ret = copy_to_user(usr_addr,
                zc info kern,
                i * sizeof(struct zc_info_elem));
```

## Appendix - compatibility

```
cmsg = CMSG_FIRSTHDR(&msg);
cmsg->cmsg_type = SCM_ZC_NOTIFICATION;

cmsg = CMSG_NXTHDR(&msg);
cmsg->cmsg_type = SO_TIMESTAMPING;

ret = sendmsg(fd, &msg, 0);
```



# Background - history of TCP zerocopy in Linux: RX

The MSG\_ZEROCOPY feature added in 4.14 enables zero-copy transmission of data, but does not address the receive side of the equation.

In linux-4.18, Eric Dumazet added a zero-copy receive mechanism to close that gap, at least for some relatively specialized applications.

- mmap() is used to reserve VMA space. tcp\_mmap() makes sure pages will be Read Only.
- getsockopt(fd, IPPROTO\_TCP, TCP\_ZEROCOPY\_RECEIVE, &zc, &zc\_len);
   To implement actual mapping of pages into user space.

## Background - history of TCP zerocopy in Linux: RX

Hardware features needed to support RX zero copy are limited to header split.

The size of the payload should be page-sized and page-aligned.