



ByteDance

RPC Latency Breaker

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ByteDance

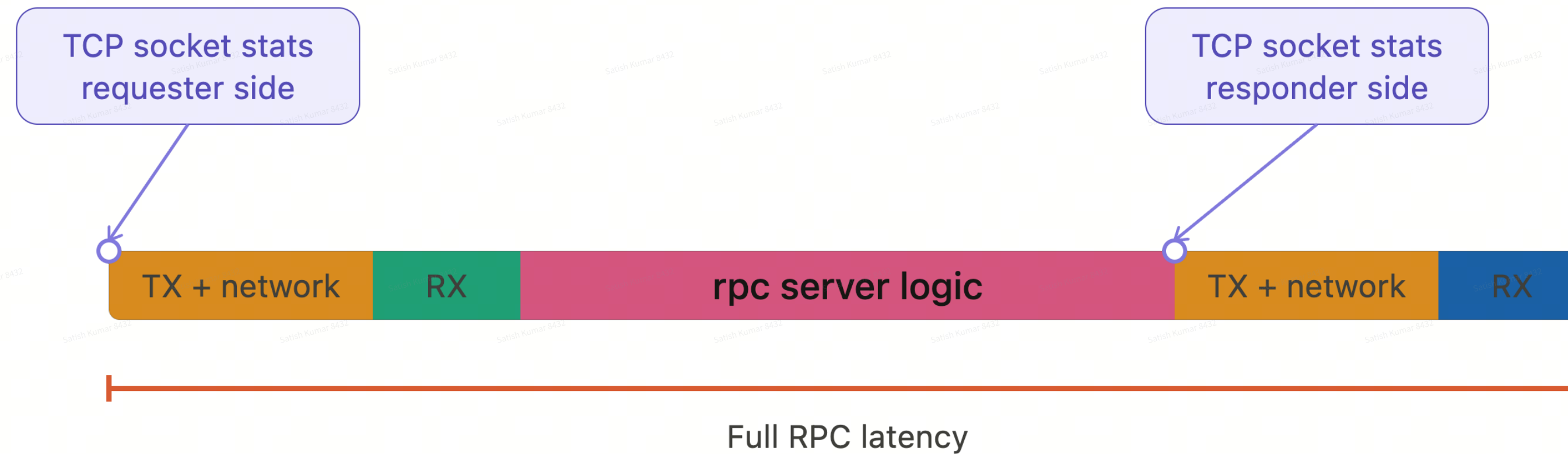
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NetDev 0x1A

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RPC Latency Breaker

Per-RPC latency breakdown



- For every RPC:
 - Full RPC latency at the system level
 - RX responder latency
 - RX requester latency
 - RPC logic latency
 - Network RTT which includes TX time
 - TCP socket stats at requester and responder



why do we care?

- RPC (remote procedure call) are fundamental communication in datacenter
 - Involve multiple servers: requester and responder
 - and multiple subsystems: network stack, scheduler, nic, network etc
- diagnosing the issue is hard:
 - multiple components involved
 - sheer amount of rpc traffic



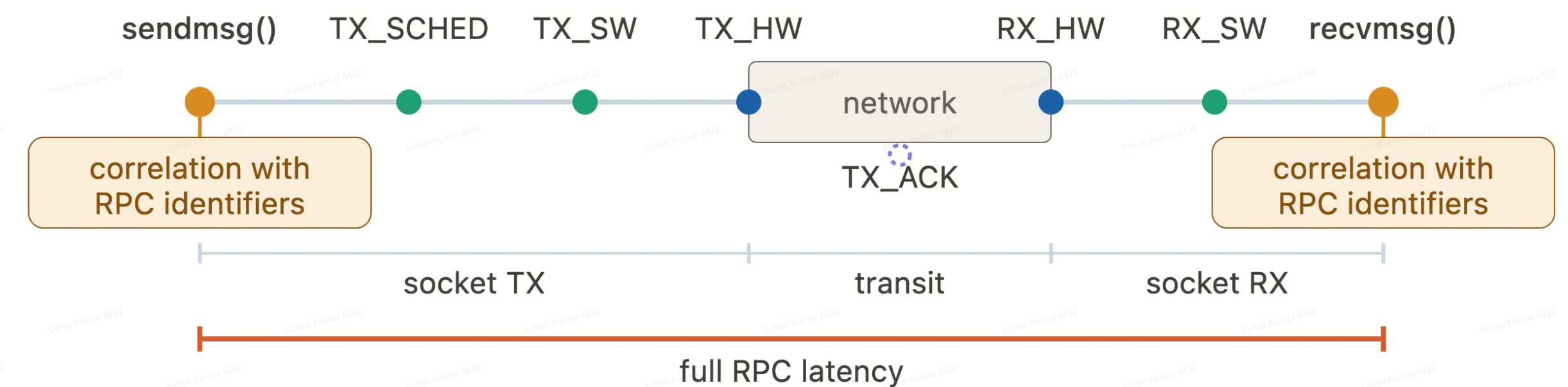
latency is the key metric

- ideally in the production
 - best average performamnce
 - no long tail latencies
- isolate the source of the issue.
 - latency breakdown
 - tcp stats

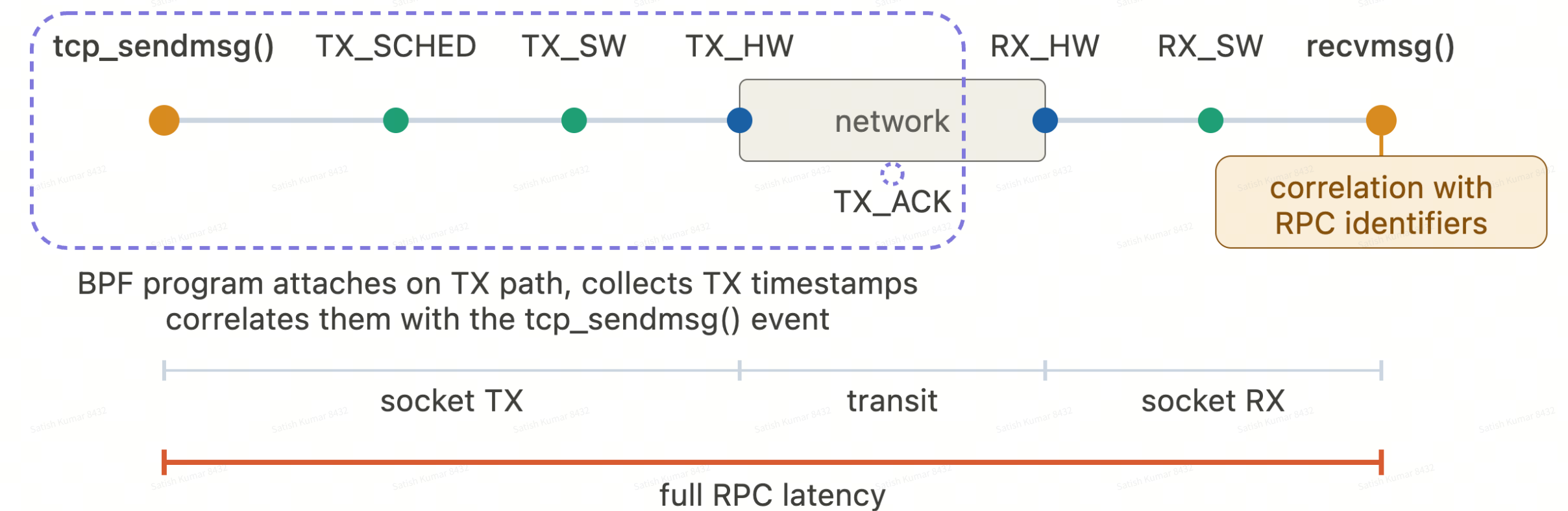
socket timestamping

- `so_timestamping`:
 - application can get various RX and TX timestamps
 - correlate with `sendmsg` and `recvmsg`
 - further correlate `sendmsg` and `recvmsg` with rpc event
 - require application changes
- `bpf so_timestamping`:
 - more efficient tx timestamp retrieval
 - can correlate timestamps with `sendmsg` event
 - cannot create rpc picture
- Reference:
 - [SO_TIMESTAMPING](#) [Willem de Bruijn]
 - [Future of SO_TIMESTAMPING](#) [Jason Xing]

SO_TIMESTAMPING → RPC latency



BPF SO_TIMESTAMPING (TX path) → RPC latency

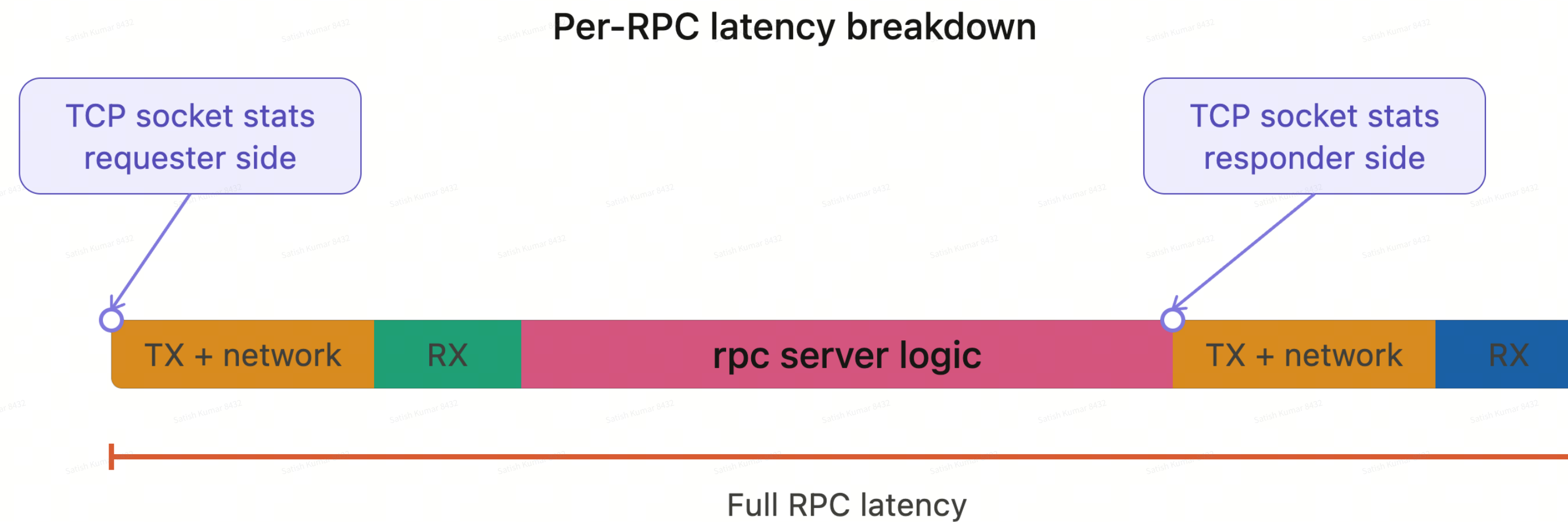




socket timestamping

- `so_timestamping`:
 - slow (tx timestamp collection is overhead)
 - application changes
- `bpf so_timestamping`:
 - fast but can not correlate with rpc

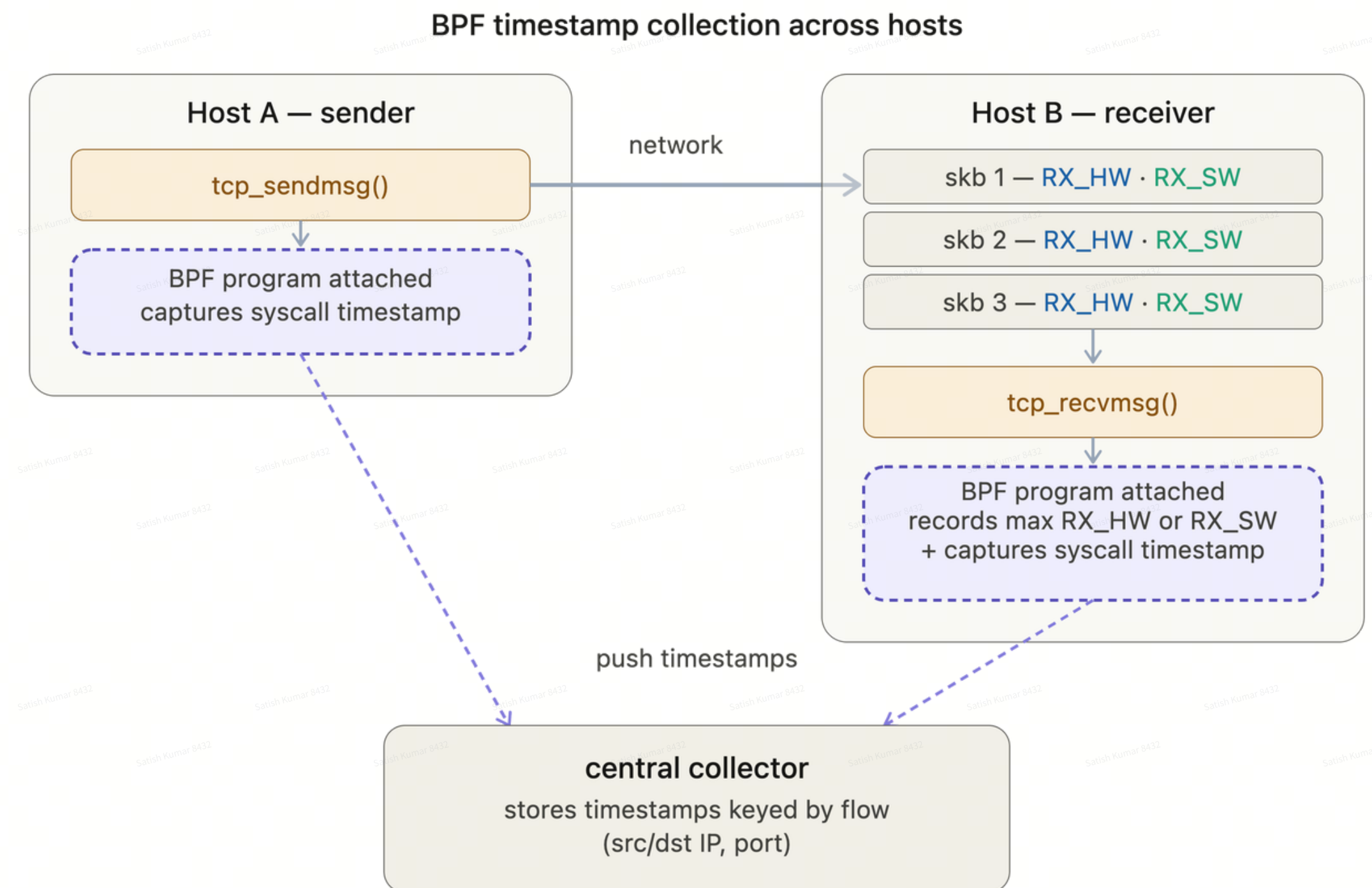
Non-intrusive, Efficient and Transparent



- rpc level correlation:
 - without application modifications -----> non-intrusive
 - efficient
 - easy to apply in production
 - cover all rpcs, no sampling
 - works for most rpc types -----> transparent

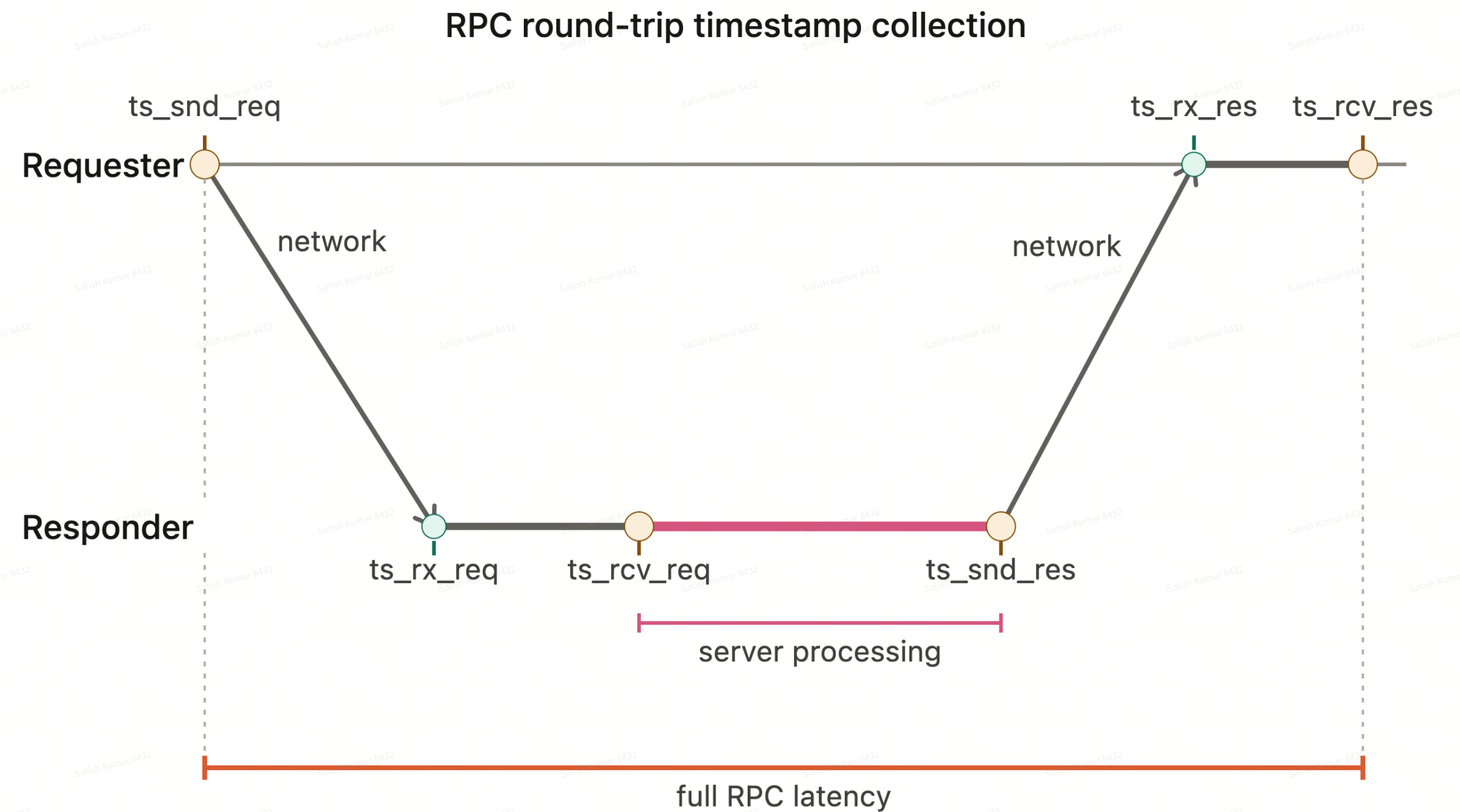
bpf program to collect timestamps

- bpf program
 - tcp_sendmsg level:
 - system call timestamp
 - tcp_recvmsg level:
 - system call timestamp
 - max of RX_HW/RX_SW so_timestamp from processed skbs



rpc timestamps

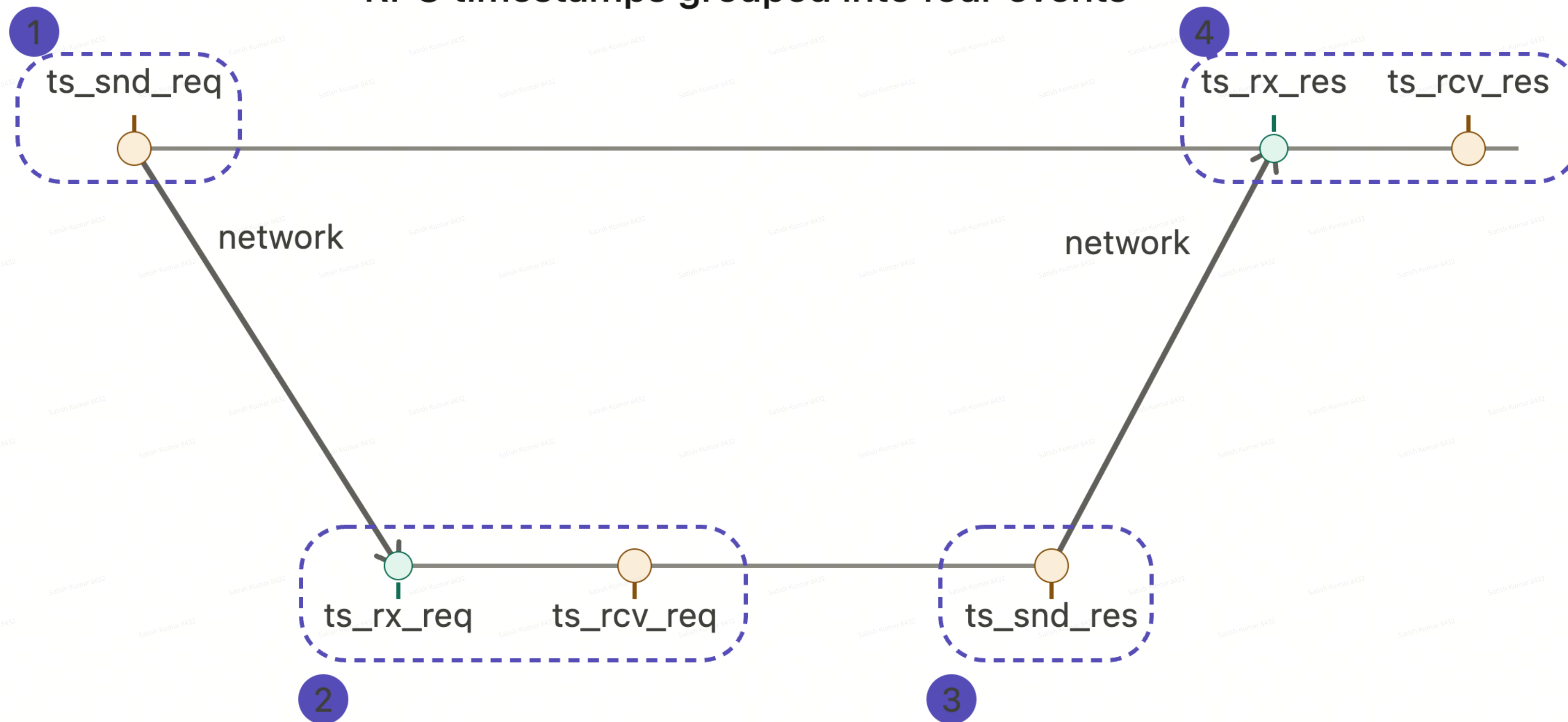
- total six timestamps per rpc
- rpc request:
 - `ts_snd_req` --> `ts_rx_req` --> `ts_rcv_req`
- rpc processing
- rpc response:
 - `ts_snd_res` --> `ts_rx_res` --> `ts_rcv_res`



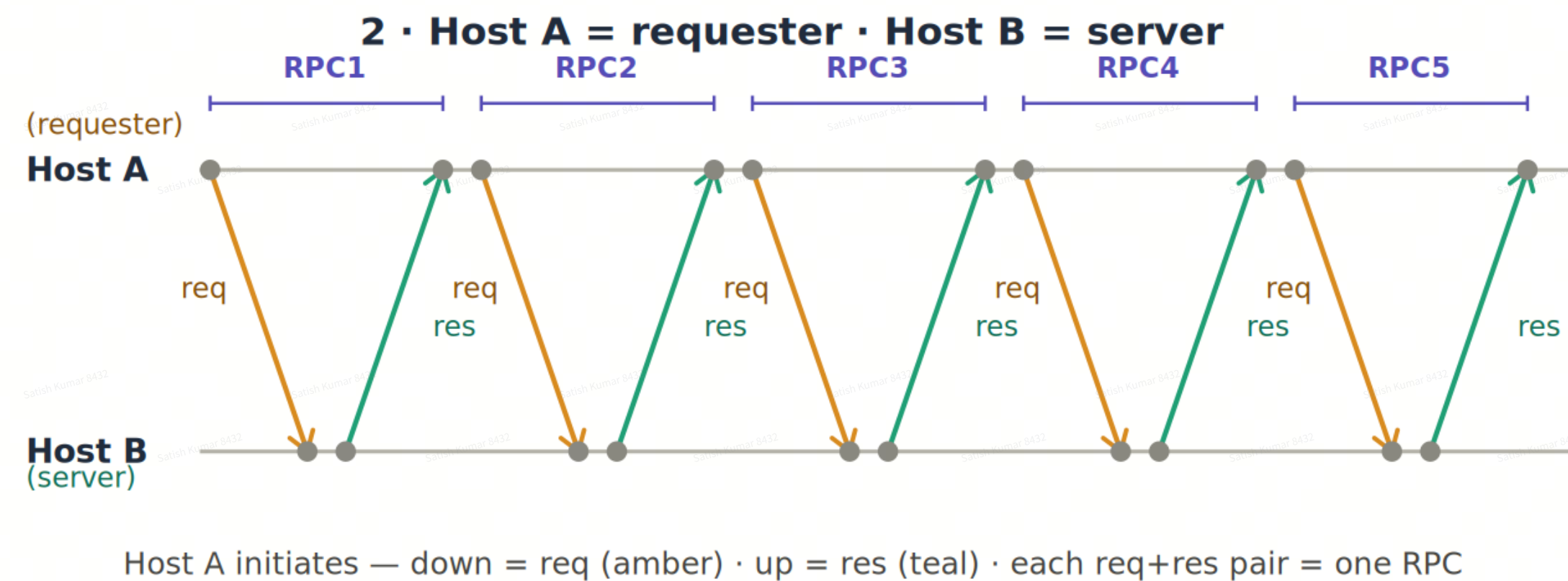
`snd` = `sendmsg()` · `rx` = received at NIC/kernel · `rcv` = read via `recvmsg()`

independent events, how to correlate?

RPC timestamps grouped into four events



rpc's ping pong pattern for correlation

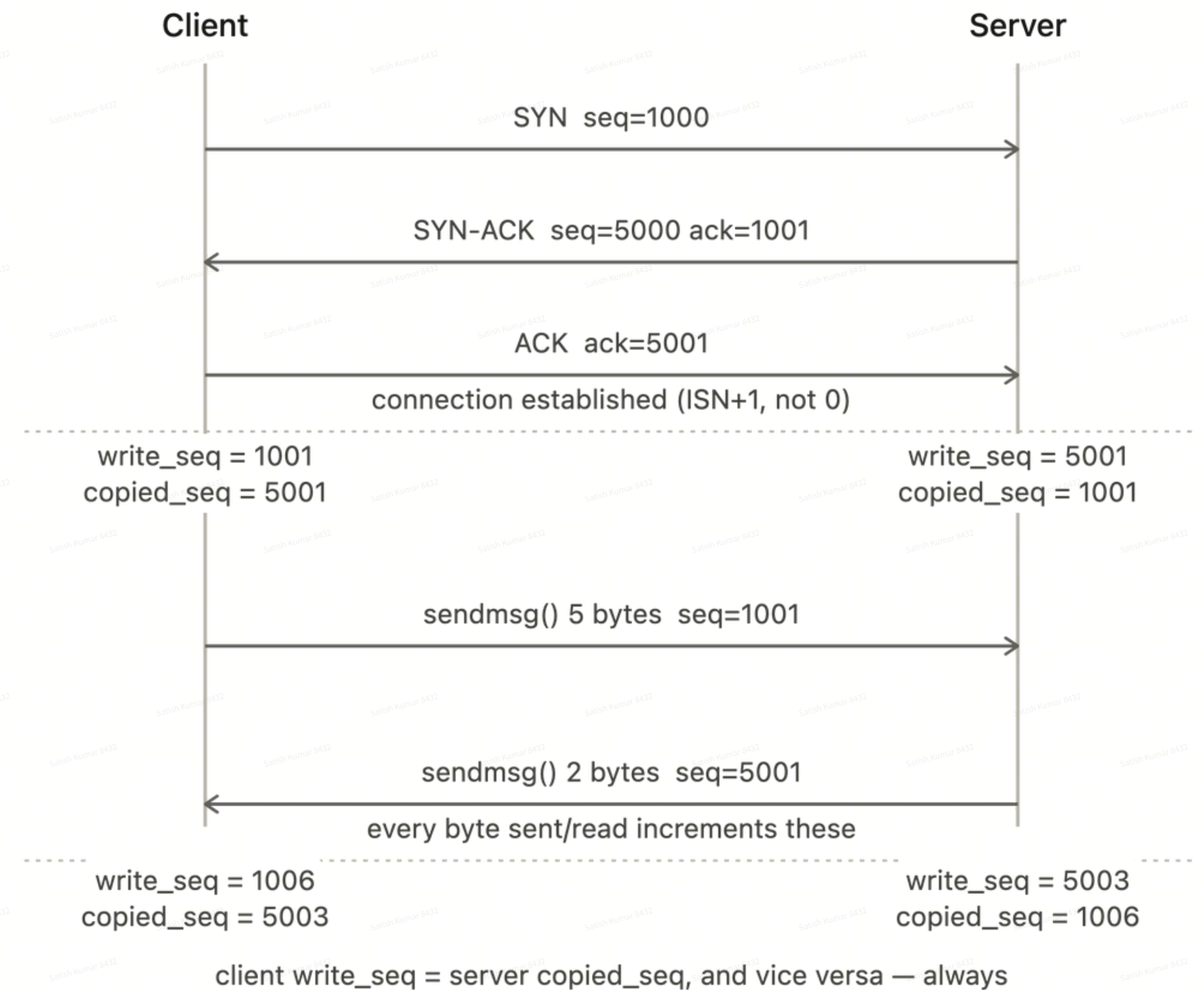


- ping: rpc request
- pong: rpc response
- its not a stream of messages towards one direction
- both side alternate between request and reponse

correlation by socket sequence number

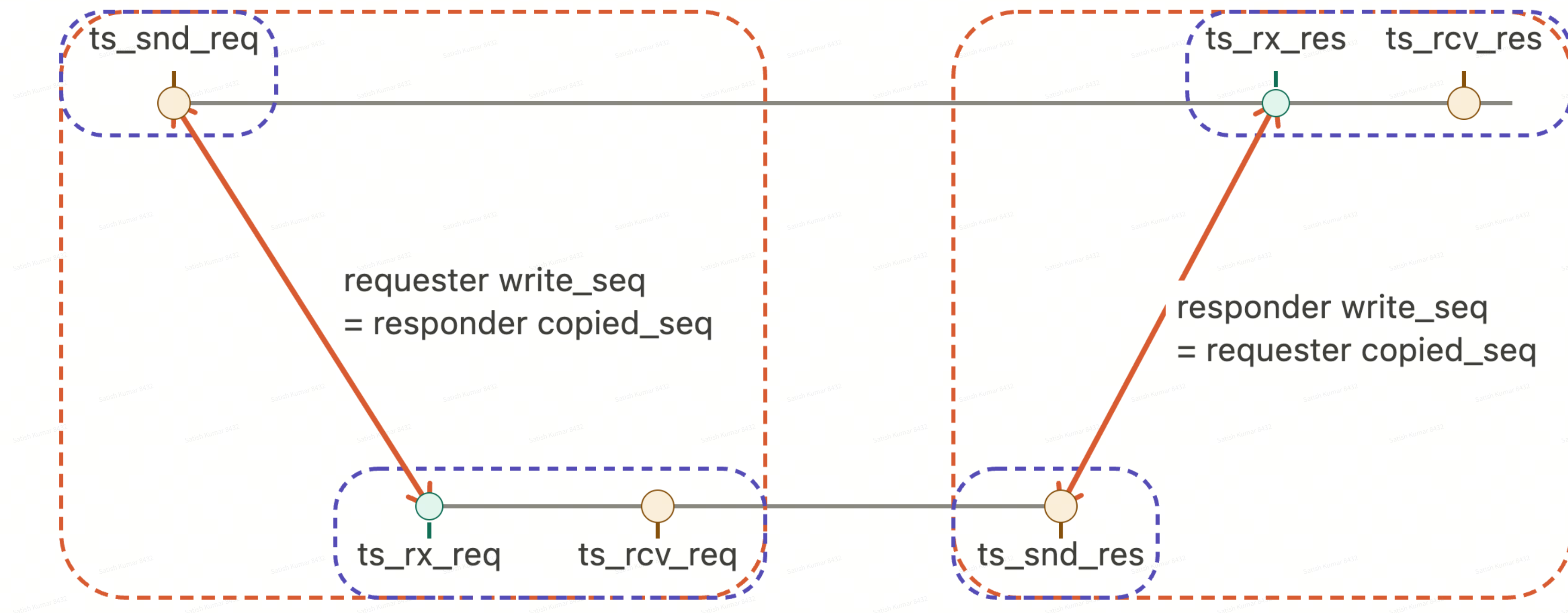
- socket sequence numbers
 - write_seq: incremented for every tcp_sendmsg byte sent by the application
 - copied_seq: incremented for every tcp_recvmsg byte received by the application
- increased when app finishes the system call
- client and server both maintain them
- and initialized during connect
 - client write_seq = server copied_seq
 - server write_seq = client copied_seq

How write_seq and copied_seq are initialized and incremented



correlate sendmsg to recvmsg events

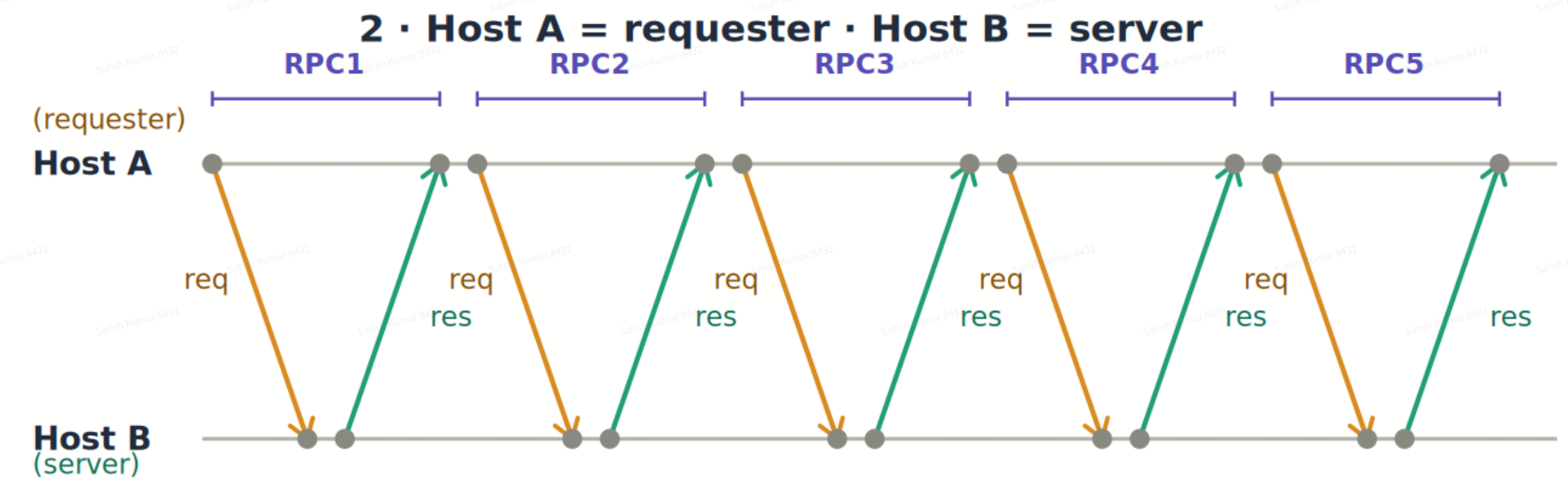
write_seq and copied_seq correlate the legs



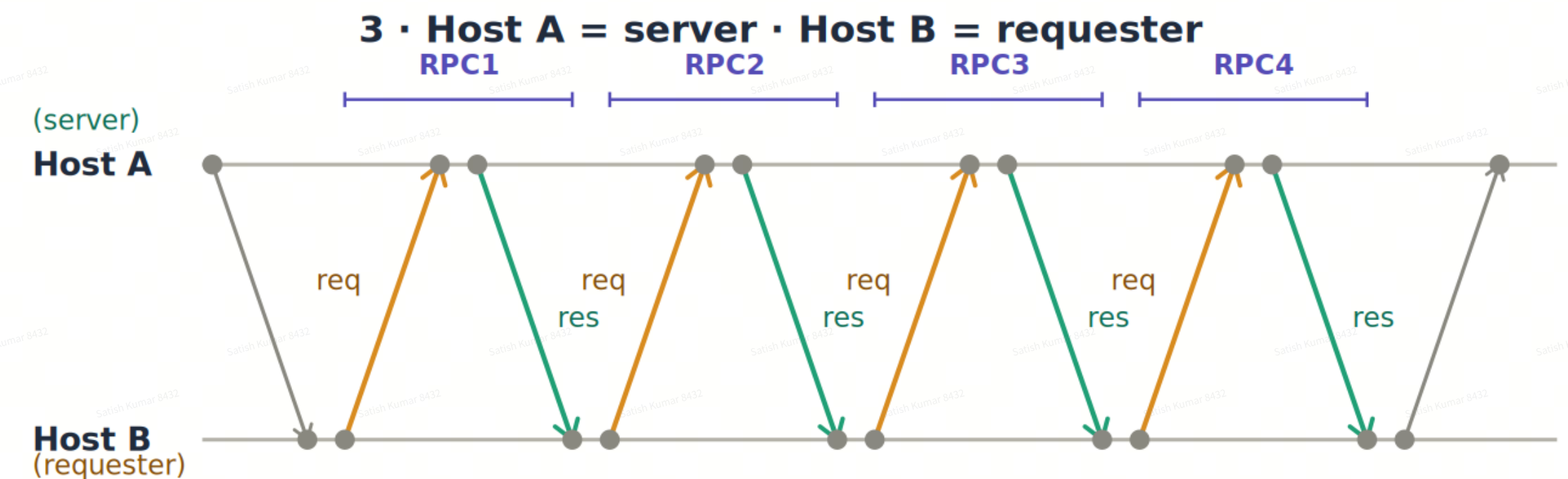
- rpc communication is sequential not stream
 - rpc request ----> rpc response
- record these sequence numbers along with timestamps

correlate request and response messages

- its just about identifying who is the rpc requester
- given ping pong behaviour, req and res can be combined to form the rpc event



Host A initiates — down = req (amber) · up = res (teal) · each req+res pair = one RPC



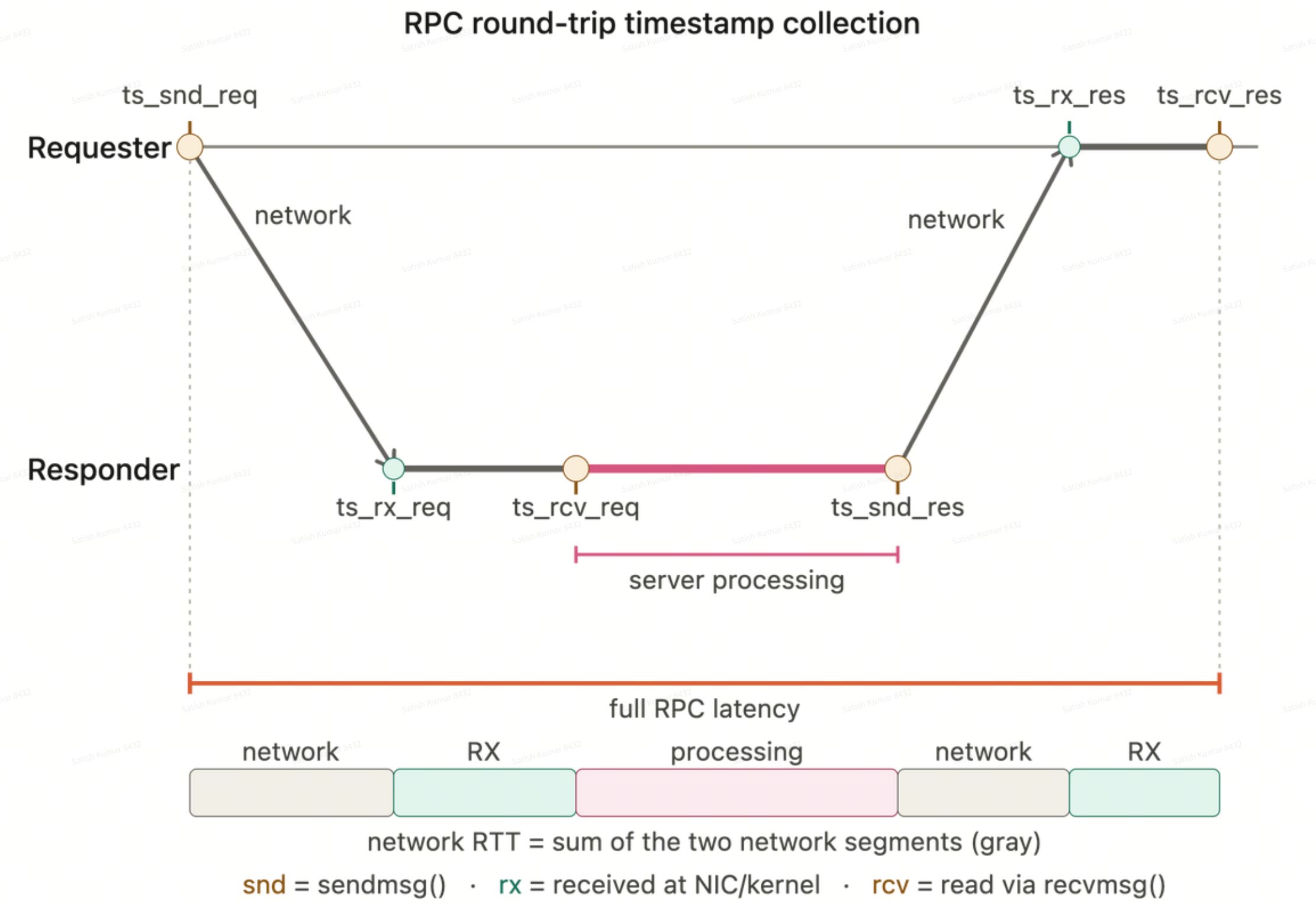
Host B initiates — up = req (amber) · down = res (teal) · gray = partial view at edges

classifying requester and responder

- we know upfront, which app is rpc requester and communicating downstream to which app

complete breakdown

- rpc latency at system level
 - $ts_rcv_res - ts_snd_req$
- rpc server processing
 - $ts_snd_res - ts_rcv_req$
- rx latencies
 - time taken since complete data arrives and application reads
 - $ts_rcv_req - ts_rx_req$
 - $ts_rcv_res - ts_rx_res$
- network rtt including tx time
 - subtract all latencies above from the total latency
 - includes the physical network delays and tx stacks





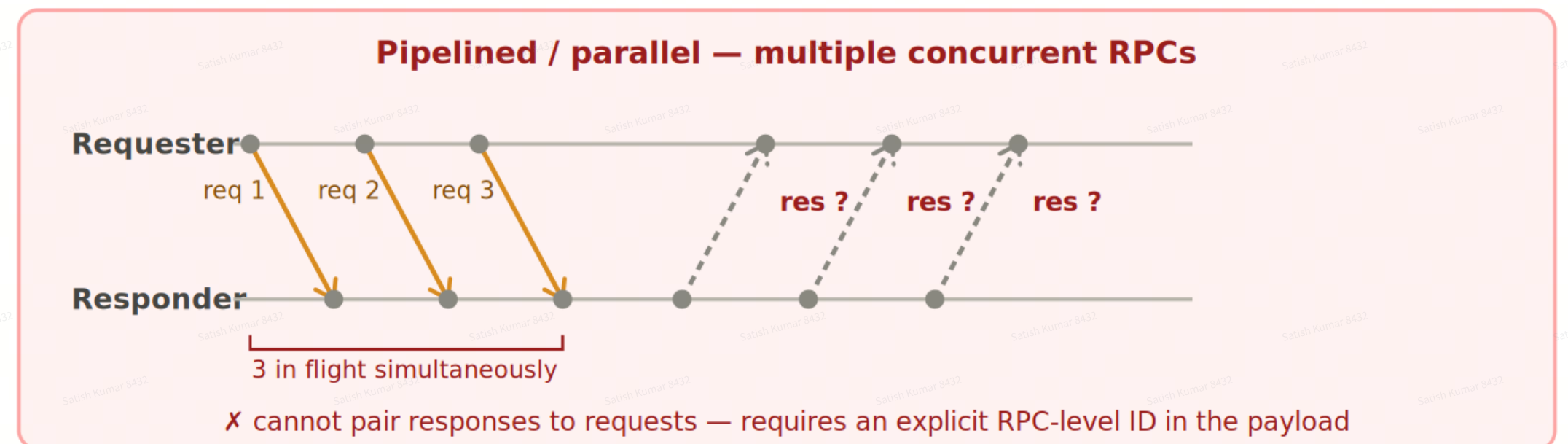
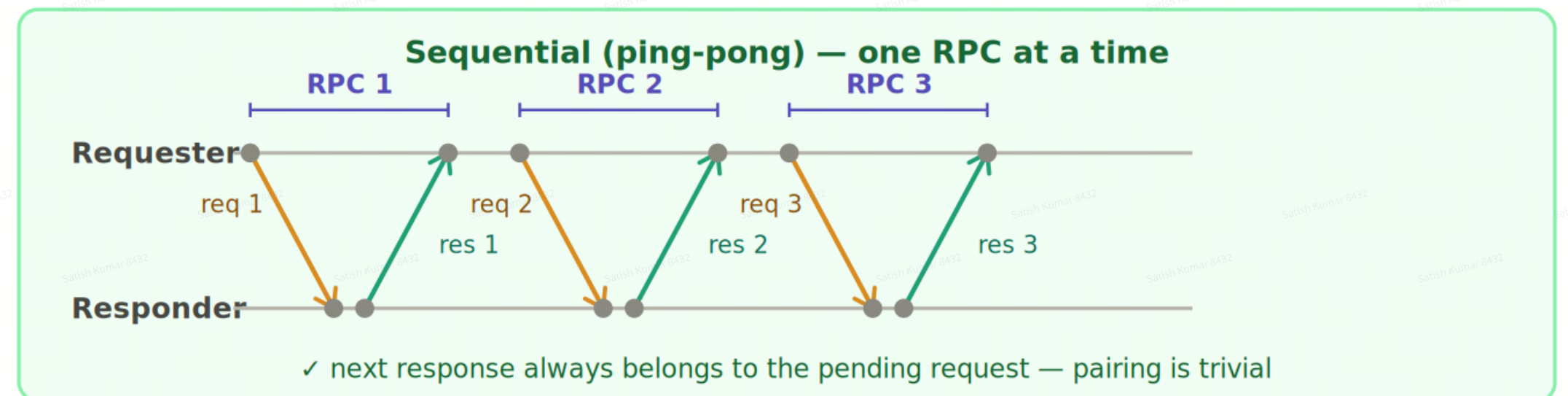
points to mention

- network rtt include tx stack
 - we could also do tx so_timestamps using bpf interface
 - we decide not to do in general due the need of packet level tracing
 - usually tx is straight path to the nic in the app context
- rpc messages are just not single sendmsg or recvmsg
 - require intelligent aggregation

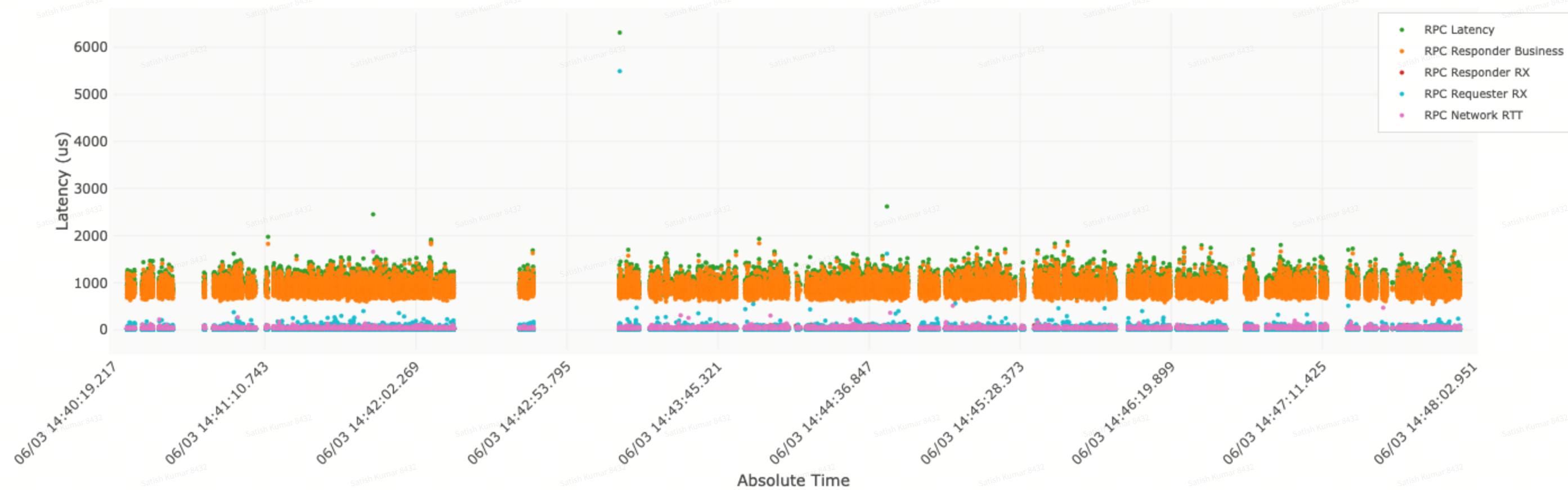
ping-pong condition

- rpc correlation requires one active call per flow
 - wont work for multiple concurrent rpc
- how common is ping pong pattern?
 - thrift protocol over tcp socket is strictly sequential by default
 - maintain connection pool
 - http/1.1 REST is also sequential and maintain connection pool
 - heavily used for internal datacenter traffic
 - grpc, http/2 REST send multiple rpc over single connection
 - given scale out container services, ping pong pattern is still common

RPC correlation requires one active call per flow

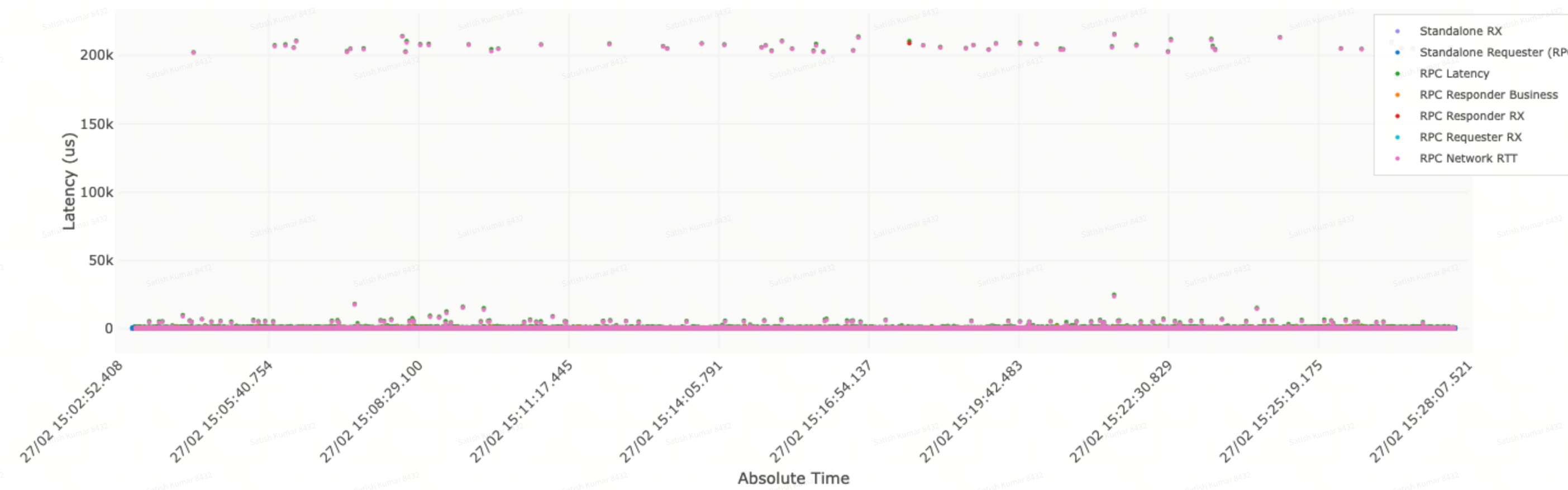


case study 1



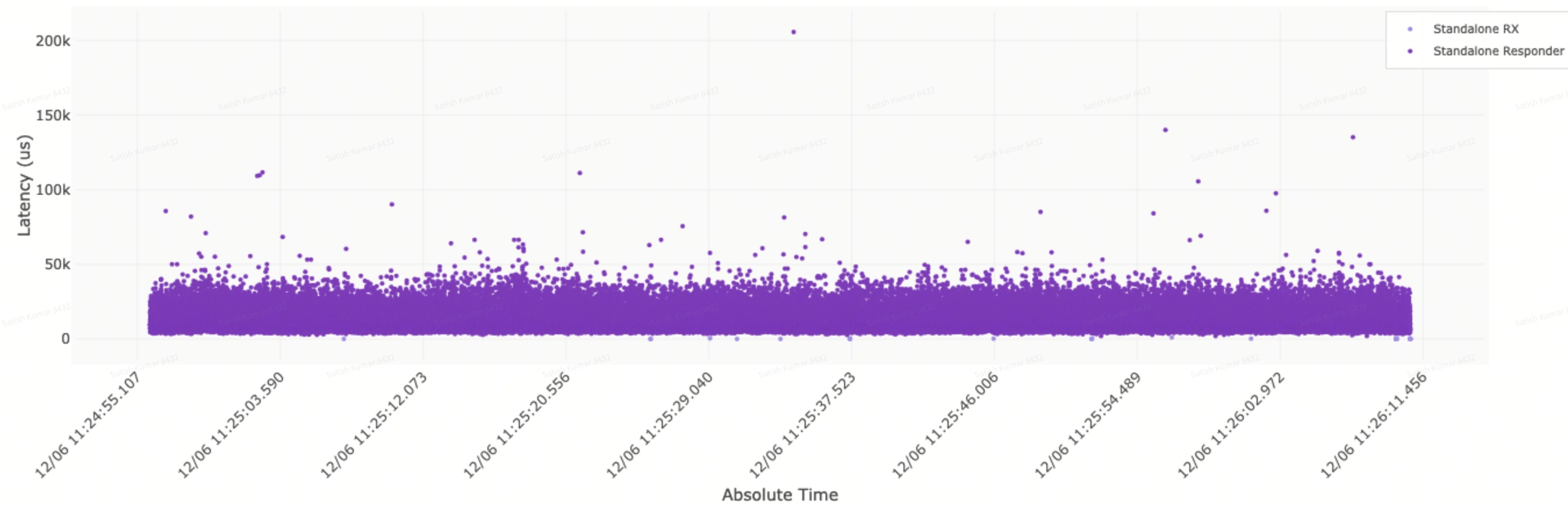
- increased requester-rx latency
 - either issue with softirq processing (rx stack) or application scheduling delay
- bcc runqlat tool and atop confirmed the scheduling delays

case study 2



- frequent jump of 200ms in network rtt
 - issue with physical network, tx stack or tcp flow
- check tcp flow stats exact at these points
 - increase in tail loss probes

case study 3



- tracing one end provide details on its communication to all other ends
- checking tcp socket stats, revealed two problem:
 - packet losses, increase in retrans
 - high tcp rtt to some machines, hardware clock issues



Thank you

Q&A